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# Negative solutions to Post's problem, II\*

By SY D. FRIEDMAN

This paper is an application of the techniques of modern set theory to problems in ordinal recursion theory. We concentrate on the global structure of the  $\beta$ -degrees for both admissible and inadmissible  $\beta$  and we show:

**THEOREM 4** ( $V = L$ ). *Let  $\beta = \aleph_{\omega_1}$  and  $0'$  = the  $\beta$ -degree of the complete  $\beta$ -r.e. set. Then the  $\beta$ -degrees greater than or equal to  $0'$  are well-ordered by  $\leq_\beta$  with successor given by the  $\beta$ -jump.*

**THEOREM 10** ( $V = L$ ). *Let  $\beta = \aleph_{\omega_1} \cdot \omega$ . There is a well-ordered sequence  $e_0 < e_1 < \dots$  of  $\beta$ -degrees such that if  $e$  is an arbitrary  $\beta$ -degree then for some  $\gamma$ ,  $e_\gamma \leq e < e_{\gamma+1}$ .*

Moreover, if we assume the Generalized Continuum Hypothesis, the  $\aleph_{\omega_1}$ -degrees and the Turing degrees are not elementarily equivalent as partial orderings. If  $V = L$  then the  $\aleph_{\omega_1}$ -jump is definable just in terms of the ordering of  $\aleph_{\omega_1}$ -degrees.

These results are in sharp contrast with earlier ones in ordinal recursion theory (see [9]), which tend to show that the  $\beta$ -degrees have a very rich structure. Thus our work here shows that the broader point of view obtained by considering inadmissible  $\beta$  has led to a structure theory of importance even for the admissible case.

In [8], Sacks and Simpson first established a connection between Gödel techniques in the study of  $L$  and ordinal recursion theory. This paper furthers this idea by relating deep results of Jensen on the fine structure of  $L$  to the structure of the  $\beta$ -degrees. The application of methods of combinatorial set theory, initiated in [1], is also greatly extended as we make use of techniques developed by Silver in [11] where he settles the singular cardinal problem at uncountable cofinalities.

Our earlier paper [2] is not a prerequisite for understanding the wo

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reported here, though we do presume a familiarity with the basic notions of  $\beta$ -recursion theory (as described in [1]). In [2], we applied Fodor's Theorem and the theory of stationary sets to establish the  $\leq_{\omega^\beta}$ -comparability of any two  $\beta$ -r.e. sets when  $\beta = \aleph_{\omega_1} \cdot \omega$ . In Section 1 of this paper similar methods are applied to establish Theorem 4.

However, the deeper results require an excursion into the fine structure of  $L$ , where the theory of master codes enables us to identify explicitly the degrees appearing in Theorems 4 and 10. This is carried out for Theorem 4 in Section 2. Theorem 10 requires the introduction of some further recursion-theoretic ideas and is dealt with in Section 3.

There are natural versions of the above results when  $\aleph_{\omega_1}$  is replaced by any singular cardinal of uncountable cofinality. However these results do not hold when  $\aleph_{\omega_1}$  is replaced by a singular cardinal of countable cofinality, as we will discuss in [3]. We end our paper by listing some further results and open questions concerning both the countable cofinality case and the finer structure of the  $\beta$ -degrees.

### 1. Stationary sets and the $\aleph_{\omega_1}$ -degrees

We assume  $V = L$ . Our first goal is to prove that if  $\alpha = \aleph_{\omega_1}$  then the  $\alpha$ -degrees  $\geq 0'$  are well-ordered by  $\leq_\alpha$ , where  $0' = \alpha$ -degree of the complete  $\alpha$ -r.e. set. This provides an example of a first-order difference between the partial-orderings of  $\aleph_{\omega_1}$ -degrees and Turing degrees; indeed, the following sentence holds in the latter but not the former:

$$\forall d \exists e \exists f (e \not\leq d \cup f \text{ and } f \not\leq d \cup e).$$

First-order differences in the language with jump were already discovered by Richard Shore (see [9]).

The proof technique used here is in fact imbedded in our earlier paper [2]. We continue to let  $\alpha$  denote  $\aleph_{\omega_1}$  (and to assume  $V = L$ ).

LEMMA 1.  $C = \{\aleph_\beta \mid \beta < \omega_1\}$  has  $\alpha$ -degree  $0'$ .

*Proof.* Note that  $C$  is  $\Pi_1$  over

$$L_\alpha: \gamma \in C \iff \gamma \geq \omega \wedge \sim \exists \gamma' < \gamma \exists f (f: \gamma \xrightarrow{1-1} \gamma').$$

So  $C \leq_\alpha 0'$ . But every element of  $C$  is  $\alpha$ -stable; i.e.,  $\gamma \in C \rightarrow L_\gamma$  is a  $\Sigma_1$ -elementary substructure of  $L_\alpha$ . (This is proved in [8]). Thus, if  $\phi(x)$  is a  $\Sigma_1$  formula defining (over  $L_\alpha$ ) a complete  $\alpha$ -r.e. set  $A$  then for  $K, H \in L_\alpha$ :

$$\begin{aligned} K \subseteq A &\leftrightarrow \exists \gamma \in C (K \subseteq L_\gamma \wedge \forall \delta \in K, L_\gamma \models \phi(\delta)), \\ H \subseteq L_\alpha - A &\leftrightarrow \exists \gamma \in C (H \subseteq L_\gamma \wedge \forall \delta \in H, L_\gamma \models \sim \phi(\delta)). \end{aligned}$$

So  $A \leq_\alpha C$ . □

Every  $A \subseteq \alpha$  of  $\alpha$ -degree  $\geq_\alpha 0'$  has the same  $\alpha$ -degree as an apparently much simpler set, its "cutoff" function.

*Definition.* For  $A \subseteq \alpha$  define  $f_A: \omega_\beta \rightarrow \alpha$ , the cutoff function for  $A$ , by  $f_A(\delta) = \beta$  if  $A \cap \aleph_\delta$  is the  $\beta$ th element in the canonical well-ordering  $<_{L_\alpha}$  of  $L_\alpha$ .

Thus we see that for all  $\delta$ ,  $f_A(\delta) < \aleph_{\delta+1}$ . Note that  $f_A \upharpoonright X$  for  $X$  unbounded in  $\omega_1$  determines  $f_A$  completely. The idea of comparing two sets  $A, B$  by comparing the growth-rates of  $f_A, f_B$  has its beginnings in Jack Silver's work on the Singular Cardinal Problem [11] and is made very explicit in Karel Prikry's proof [7] of Silver's theorem. It is this idea which we use here.

LEMMA 2. *Suppose  $0' \leq_\alpha A, 0' \leq_\alpha B$ . Then either  $A \leq_\alpha B$  or  $B \leq_\alpha A$ .*

*Proof.* One of the sets  $\{\delta \mid f_A(\delta) \leq f_B(\delta)\}, \{\delta \mid f_A(\delta) \geq f_B(\delta)\}$  must be a stationary subset of  $\omega_1$ . We assume that the former set is and proceed to show that  $A \leq_\alpha B$ .

It is enough to show that  $f_A \upharpoonright X$  can be computed from  $B$  for some unbounded  $X \subseteq \omega_1$ .

For any  $\alpha < \beta$ , let  $c_\beta$  be the  $<_L$ -least injection of  $\beta$  into cardinality  $(\beta)$ . Define functions  $g, h$  on  $\{\delta \mid f_A(\delta) \leq f_B(\delta)\}$  by

$$\begin{aligned} g(\delta) &= c_{f_B(\delta)+1}(f_A(\delta)), \\ h(\delta) &= \text{least } \delta' \text{ such that } g(\delta) < \aleph_{\delta'}. \end{aligned}$$

Then  $g(\delta) < \aleph_\delta$  for all  $\delta$ , so  $\delta \text{ limit} \rightarrow h(\delta) < \delta$ . By Fodor's Theorem, choose a stationary  $X \subseteq \omega_1$  and  $\delta_0$  such that  $\delta \in X \rightarrow h(\delta) < \delta_0$ .

But then  $\delta \in X \rightarrow g(\delta) < \aleph_{\delta_0}$  so  $g \upharpoonright X$  is  $\alpha$ -finite (since any bounded subset of  $\alpha$  is a  $\alpha$ -finite). We can now compute  $f_A \upharpoonright X$  from  $B$  by

$$f_A(\delta) = c_{f_B^{-1}(\delta)+1}(g(\delta)), \quad \delta \in X. \quad \square$$

Note the heavy use of uncountable cofinality here (in order to apply Fodor's Theorem). It is known that Lemma 2 fails when  $\alpha$  is replaced by  $\aleph_\omega$ . (This is due to Leo Harrington and Bob Solovay independently and will be shown in [3].)

Our proof in fact shows:

LEMMA 3. *The  $\alpha$ -degrees  $\geq_\alpha 0'$  are well-ordered by  $\leq_\alpha$ .*

*Proof.* The proof of Lemma 2 shows that if  $0' \leq_\alpha A <_\alpha B$  then  $\{\delta \mid f_A(\delta) < f_B(\delta)\}$  contains a closed unbounded subset. Now if  $A_0 >_\alpha A >_\alpha \dots$ , each  $A_i \geq_\alpha 0'$ , then  $\{\delta \mid f_{A_{i+1}}(\delta) < f_{A_i}(\delta)\}$  contains a closed unbounded set for each  $i$  so there must be  $\delta$  such that  $f_{A_{i+1}}(\delta) < f_{A_i}(\delta)$  for every  $i$ , as the countable intersection of closed unbounded sets is nonempty. Of course we have now

contradicted the well-foundedness of the ordinals.  $\square$

It will be shown in [3] that the  $\aleph_\omega$ -degrees  $\geq 0'$  are not well-founded.

We next establish that the successor of each  $\alpha$ -degree  $\geq_\alpha 0'$  is its  $\alpha$ -jump. Our original proof of this fact made heavy use of the fine structure theory. Tony Martin later found a combinatorial proof. Subsequently we discovered a simpler combinatorial argument which we present here.

**THEOREM 4.** *The  $\alpha$ -degrees  $\geq_\alpha 0'$  are well-ordered by  $\leq_\alpha$  with successor given by the  $\alpha$ -jump.*

*Proof.* We have only to show that  $0' \leq_\alpha A <_\alpha B$  implies  $A' \leq_\alpha B$  where  $A' = \alpha$ -jump of  $A$ . Let  $\varphi(x)$  be a  $\Sigma_1$  formula which defines  $A'$  over  $\langle L_\alpha, A \rangle$ ; i.e.,  $A' = \{x \mid \langle L_\alpha, A \rangle \models \varphi(x)\}$ . For any  $\gamma < \omega_1$  we let  $A_\gamma = \{x \mid \langle L_{\aleph_\gamma}, A \cap \aleph_\gamma \rangle \models \varphi(x)\}$  and  $f_\gamma = f_{A_\gamma}$ , the cutoff function for  $A_\gamma$ . Then the sequence  $\langle A_\gamma \mid \gamma < \omega_1 \rangle$  is  $\alpha$ -recursive in  $A$ . Thus for any  $g: \omega_1 \rightarrow \omega_1$ , the function  $f_g$  defined by  $f_g(\delta) = f_{g(\delta)}(\delta)$  is also  $\alpha$ -recursive in  $A$  (as any such  $g$  is  $\alpha$ -finite).

We compare  $f_B$  to  $\sup_\gamma f_\gamma$ : If  $\{\delta \mid f_B(\delta) < \sup_\gamma f_\gamma(\delta)\}$  is stationary then for some  $g: \omega_1 \rightarrow \omega_1$  the set  $\{\delta \mid f_B(\delta) < f_g(\delta)\}$  is stationary. The proof of Lemma 2 actually shows: If  $\{\delta \mid f_B(\delta) \leq h(\delta)\}$  is stationary and  $h \leq_\alpha A$  then  $B \leq_\alpha A$ . Thus as we have assumed that  $A <_\alpha B$  it must be the case that  $\{\delta \mid \sup_\gamma f_\gamma(\delta) \leq f_B(\delta)\}$  contains a closed, unbounded set.

As in the proof of Lemma 2, for any  $\beta < \alpha$  let  $c_\beta$  denote the  $<_{\iota}$ -least injection of  $\beta$  into cardinality  $(\beta)$ . Now for limit  $\delta < \omega_1$ ,  $\aleph_\delta$  has countable cofinality. Thus for stationary many  $\delta$  we may choose an uncountable  $X_\delta \subseteq \omega_1$  and  $h(\delta) < \delta$  such that:

$$\gamma \in X_\delta \longrightarrow c_{f_{B(\delta)+1}}(f_\gamma(\delta)) < \aleph_{h(\delta)}.$$

Thus by Fodor's Theorem we may choose  $\delta_0 < \omega_1$  such that for stationary many  $\delta$ :

$$(*) \quad \gamma \in X_\delta \longrightarrow c_{f_{B(\delta)+1}}(f_\gamma(\delta)) < \aleph_{h(\delta_0)}.$$

We let  $g(\delta, \gamma) = c_{f_{B(\delta)+1}}(f_\gamma(\delta))$  and  $S = \{\delta \mid (*) \text{ holds}\}$ . For each  $\beta < \alpha$  write  $K(\beta) = y$  if  $y$  is the  $\beta$ th set in  $<_{\iota}$ . Then we see that for  $\delta \in S$ :

$$A' \cap \aleph_\delta = \bigcup_{\gamma \in X_\delta} K(c_{f_{B(\delta)+1}}^{-1}(g(\delta, \gamma))).$$

As  $g, S$  and  $\langle X_\delta \mid \delta \in S \rangle$  are  $\alpha$ -finite and  $K$  is  $\alpha$ -recursive this shows that  $A' \leq_\alpha B$ .  $\square$

We note that there is an appropriate version of Theorem 4 if only the generalized continuum hypothesis (GCH) is assumed. In this case choose  $A \subseteq \alpha$  such that  $\{\aleph_\delta \mid \delta < \omega_1\} \leq_\alpha A$  and  $2^{\aleph_\delta} \subseteq L_{\aleph_{\delta+1}}[A]$  for each  $\delta < \omega_1$ . Then Theorem 4 holds when  $0'$  is replaced by  $\alpha$ -degree  $(A)$ . Thus the GCH alone

implies that there is a first-order difference between the  $\aleph_{\omega_1}$ -degrees and the Turing degrees as partial orderings.

Next we derive some information concerning the  $\alpha$ -jump, assuming  $V = L$ . Following Carl Jockusch and David Posner, define an  $\alpha$ -degree  $d$  to be *generalized low* if  $d' = d \vee 0'$ . If  $0' \leq d$  then  $d$  is *not* generalized low. However this is the only exception.

**THEOREM 5.** *For all  $d$  either  $0' \leq d$  or  $d$  is generalized low.*

*Proof.* Define  $A \subseteq \alpha$  to be *hyper-regular* if  $\langle L_\alpha, A \rangle$  is admissible. If  $A$  is hyper-regular then  $B$  is *A-hyper-regular* if  $A \vee B$  is hyper-regular. The following claim generalizes a result of Richard Shore:

**CLAIM.** *If  $A$  is hyper-regular and  $B$  is not A-hyper-regular then  $A' \leq_\alpha B \vee A$ .*

*Proof of Claim.* Choose a function  $f: \gamma_0 \rightarrow \alpha$ ,  $\gamma_0 < \alpha$ , such that  $f$  is weakly  $\alpha$ -recursive in  $A \vee B$  and  $f$  has unbounded range. Also let  $\varphi(x)$  be a  $\Sigma_1$  formula such that  $A' = \{x \mid \langle L_\alpha, A \rangle \models \varphi(x)\}$ . As  $A$  is hyper-regular we can choose a function  $h: \gamma_0 \rightarrow \gamma_0$  such that for all  $\gamma < \gamma_0$ ,

$$A' \cap f(\gamma) = \{x \mid \langle L_{f(h(\gamma))}, A \cap f(h(\gamma)) \rangle \models \varphi(x)\}.$$

As  $h$  is  $\alpha$ -finite this shows that  $A' \leq_\alpha A \vee B$ . This proves the claim.

Now choose a set  $A$  in the  $\alpha$ -degree  $d$ . If  $A$  is not hyper-regular then  $0' \leq_\alpha A$  by the claim. Otherwise  $A' \leq_\alpha 0' \vee A$  by the claim since  $0'$  is not  $A$ -hyper-regular (Lemma 1).  $\square$

Stephen Simpson's work in [13] completes the picture by showing that every  $d \geq 0'$  is the  $\alpha$ -jump of some generalized low degree. Now a theorem of Richard Shore in [10] shows that if  $A$  is hyper-regular then there are sets  $B, C$  which are  $\alpha$ -r.e. in  $A$  such that  $B \not\leq_\alpha A \vee C$ ,  $C \not\leq_\alpha A \vee B$ . These facts enable us to prove the following theorem:

**THEOREM 6.** *The  $\alpha$ -jump operation on the  $\alpha$ -degrees is first-order definable over the structure  $\langle \alpha\text{-degrees}, \leq_\alpha \rangle$ .*

*Proof.* By our above remarks and Lemma 2 we have:

$$0' = \text{least } d \text{ such that } \forall e \forall f (d \vee e \leq_\alpha d \vee f \text{ or } d \vee f \leq_\alpha d \vee e).$$

Thus  $\{0'\}$  is first-order definable over  $\langle \alpha\text{-degrees}, \leq_\alpha \rangle$ . But then by Theorem 5:

$$e = d' \leftrightarrow [0' \not\leq_\alpha d \text{ and } e = d \vee 0'] \text{ or } [0' \leq_\alpha d \text{ and } e = \text{least degree } > d]. \quad \square$$

We end this section by briefly mentioning how the above results extend

to other singular cardinals of uncountable cofinality. Let  $\beta$  be such a cardinal. Then there is a least  $\beta$ -degree  $d$  such that the  $\beta$ -degrees  $\geq d$  are well-ordered by  $\leq_\beta$  (the least non-hyper-regular  $\beta$ -degree). Moreover for  $A \subseteq \beta$  one can make an appropriate definition of  $A^{(\nu)}$ , the  $\nu$ th iterate of the  $\beta$ -jump applied to  $A$ , and then  $d = 0^{(\nu)}$  for some  $\nu < \beta^+ = \text{next cardinal after } \beta$ . If  $d \not\leq_\beta A$  then  $A$  is generalized low $_\nu$  (i.e.,  $A^{(\nu)} =_\beta A \vee 0^{(\nu)}$ ) and hyper-regular.  $\{d\}$  is definable over  $\langle \beta\text{-degrees, } \leq_\beta \rangle$ .

## 2. Master codes

Let  $\beta$  be a limit ordinal. We define the  $\Sigma_n$  *projectum* of  $\beta$ ,  $\rho_n^\beta$ , and the  $\Delta_n$  *projectum* of  $\beta$ ,  $\delta_n^\beta$ , by

$\rho_n^\beta = \text{least } \gamma \text{ such that there is a 1-1 function from } \beta \text{ into } \gamma$   
which is  $\Sigma_n$  over  $S_\beta$

$\delta_n^\beta = \text{least } \gamma \text{ such that there is a 1-1 function from } \beta \text{ onto } \gamma$   
which is  $\Sigma_n$  over  $S_\beta$ .

We also let  $\rho_0^\beta = \delta_0^\beta = \beta$ .

The  $S$ -hierarchy for  $L$  is defined in Devlin's book [1], page 82. It is often easier to work with this hierarchy as  $S_\beta \cap \text{On} = \beta$  for limit  $\beta$ . (For all  $\beta$ ,  $S_{\omega_\beta} = J_\beta$  so the  $S$ -hierarchy is merely a "ramification" of the  $J$ -hierarchy.) Our definition of the projecta differs from Jensen's but is easily related to his definition and allows one to state cleanly the following characterization, which follows from Jensen's work:

**THEOREM 7.** (a)  $\rho_n^\beta = \text{least } \gamma \text{ such that there is a subset of } \gamma \text{ which is } \Sigma_n \text{ over } S_\beta \text{ but not a member of } S_\beta$ .

(b)  $\delta_n^\beta = \text{least } \gamma \text{ such that there is a subset of } \gamma \text{ which is } \Delta_n \text{ over } S_\beta \text{ but not a member of } S_\beta$ .

(c) For  $\delta \leq \beta$  let  $\Sigma_n^\beta \text{ cf}(\delta) = \text{least } \gamma \text{ such that there is a function from } \gamma \text{ onto an unbounded subset of } \delta \text{ which is } \Sigma_n \text{ over } S_\beta$ . Then  $\delta_n^\beta = \max(\rho_n^\beta, \Sigma_n^\beta \text{ cf}(\rho_{n-1}^\beta))$ .

The notion of a  $\Sigma_n$  master code is the key in Jensen's fine structure theory.  $A \subseteq \beta$  is a  $\Sigma_n$  master code for  $\beta$  if:

(a)  $A \subseteq \rho_n^\beta$ ; (b)  $A$  is  $\Sigma_n$  over  $S_\beta$ ; (c) Let  $\mathfrak{A} = \langle S_{\rho_n^\beta}, A \rangle$ . Then  $B \subseteq \rho_n^\beta$  is  $\Sigma_1$  over  $\mathfrak{A}$  if and only if  $B$  is  $\Sigma_{n+1}$  over  $S_\beta$ .

Jensen showed that  $\Sigma_n$  master codes always exist. As in [4], choose a canonical  $\Sigma_n$  master code for  $\beta$ ,  $C_n^\beta$ , and let  $\mathfrak{A}_n^\beta = \langle S_{\rho_n^\beta}, C_n^\beta \rangle$ . It follows from (c) that  $B \subseteq \rho_n^\beta$  is  $\Sigma_m$  over  $\mathfrak{A}_n^\beta$  if and only if  $B$  is  $\Sigma_{n+m}$  over  $S_\beta$ .

The notion of a  $\Delta_n$  master code occurs more rarely (though it was

considered by Jockusch and Simpson [5] in case  $\rho_n^\beta = \omega$ ).  $A \subseteq \beta$  is a  $\Delta_n$  master code for  $\beta$  if:

(a)  $A \subseteq \delta_n^\beta$ , (b)  $A$  is  $\Delta_n$  over  $S_\beta$ . (c) If  $\mathfrak{A} = \langle S_{\delta_n^\beta}, A \rangle$  then  $B \subseteq \delta_n^\beta$  is  $\Sigma_1$  over  $\mathfrak{A}$  if and only if  $B$  is  $\Sigma_n$  over  $S_\beta$ .

Again, (c) implies that  $B \subseteq \delta_n^\beta$  is  $\Sigma_m$  over  $\mathfrak{A}$  if and only if  $B$  is  $\Sigma_{m+n-1}$  over  $S_\beta$ .  $\Delta_n$  master codes do not always exist. They do if  $\delta_n^\beta = \omega$  or if  $\rho_{n-1}^\beta = \delta_n^\beta$  (in this case  $C_{n-1}^\beta$  is a  $\Delta_n$  master code for  $\beta$ ). The case that concerns us here is when  $\delta_n^\beta = \aleph_{\omega_1}$  and the results below show that a  $\Delta_n$  master code exists for  $\beta$  if and only if  $\Sigma_n^{\beta}cf(\rho_{n-1}^\beta) = \omega_1$ . In general a  $\Delta_n$  master code exists for  $\beta$  if and only if  $\Sigma_n^{\beta}cf(\rho_{n-1}^\beta) = \Sigma_n^{\beta}cf(\delta_n^\beta)$ .

As in the previous section let  $\alpha$  denote  $\aleph_{\omega_1}$  and assume  $V = L$ . We will describe a procedure for dissecting the  $\alpha$ -degrees  $\geq 0'$ . We first establish some basic facts about master codes which are subsets of  $\alpha$ ,  $\alpha$ -master codes. Define  $A$  to be an  $\alpha$ -master code for  $\beta$  if  $A$  is a  $\Sigma_n$  or  $\Delta_n$  master code for  $\beta$  for some  $n$  and  $A \subseteq \alpha$ .

**PROPOSITION 8.** (a) If  $\rho_n^\beta = \alpha$  then all  $\Sigma_n$  master codes for  $\beta$  have the same  $\alpha$ -degree.

(b) If  $\delta_n^\beta = \alpha$  then all  $\Delta_n$  master codes for  $\beta$  have the same  $\alpha$ -degree.

(c) If  $\beta_1 < \beta_2$  then all  $\alpha$ -master codes for  $\beta_1$  are  $\alpha$ -recursive in all  $\alpha$ -master codes for  $\beta_2$ .

(d) If  $\rho_{n-1}^\beta = \alpha$  then the  $\Sigma_{n-1}$  master codes for  $\beta$  and the  $\Delta_n$  master codes for  $\beta$  have the same  $\alpha$ -degree.

(e) If  $\alpha = \delta_n^\beta$  and if  $D$  is a  $\Delta_n$  master code for  $\beta$ ,  $C$  a  $\Sigma_n$  master code for  $\beta$ , then  $C = {}_\alpha \alpha$ -jump( $D$ ).

(f) If  $\rho_n^\beta = \alpha$  and  $C_1$  is a  $\Sigma_n$  master code for  $\beta$ ,  $C_2$  a  $\Sigma_{n+1}$  master code for  $\beta$ , then  $C_2 = {}_\alpha \alpha$ -jump( $C_1$ ).

Proposition 8 gives us the following picture: Say that a limit ordinal  $\beta \geq \alpha$  is projectible into  $\alpha$  if  $\rho_n^\beta = \alpha$  for some  $n$ . List the limit ordinals  $\beta \geq \alpha$  which are projectible into  $\alpha$  in order:  $\alpha = \beta_0 < \beta_1 < \dots$  and list the  $\alpha$ -degrees of  $\alpha$ -master codes in  $\leq_\alpha$ -increasing order  $0 = d_1 < d_2 < \dots$ . If  $\beta$  is projectible into  $\alpha$  let  $n(\beta) =$  least  $n$  such that  $\rho_n^\beta = \alpha$ . Now let  $\beta = \beta_\nu$ ,  $n = n(\beta)$  and  $\lambda = \omega \cdot \nu$ . Then if  $\delta_n^\beta = \alpha$  and there is a  $\Delta_n$  master code for  $\beta$ , we have  $d_\lambda = \alpha$ -degree (any  $\Delta_n$  master code for  $\beta$ ) and  $d_{\lambda+m} = \alpha$ -degree (any  $\Delta_{n+m}$  master code for  $\beta$ ). This occurs exactly when  $\Sigma_n^{\beta}cf(\rho_{n-1}^\beta) = \omega_1$ . Otherwise  $d_{\lambda+m} = \alpha$ -degree (any  $\Sigma_{n+m}$  master code for  $\beta$ ). Finally, for all  $\nu$ ,  $d_{\nu+1} = \alpha$ -jump( $d_\nu$ ).

Our ultimate aim is to show that the  $d_\nu$ 's,  $\nu \geq 2$ , exhaust the  $\alpha$ -degrees  $\geq {}_\alpha 0' = d_2$ . We show how this is done after proving Proposition 8.



*Proof of Proposition 8.* (a) If  $A \subseteq L_\alpha$  then  $\{x \in L_\alpha \mid x \subseteq A\}$  and  $\{x \in L_\alpha \mid x \cap A = \emptyset\}$  are both  $\Sigma_1$  over  $\langle L_\alpha, A \rangle$ . Therefore if  $C_1, C_2$  are  $\Sigma_n$  master codes for  $\beta$ ,  $\rho_n^\beta = \alpha$  then  $\{x \in L_\alpha \mid x \subseteq C_1\}$  and  $\{x \in L_\alpha \mid x \cap C_1 = \emptyset\}$  are  $\Sigma_1$  over  $\langle L_\alpha, C_1 \rangle$ , hence  $\Sigma_{n+1}$  over  $\beta$ , hence  $\Sigma_1$  over  $\langle L_\alpha, C_2 \rangle$ . So  $C_1 \leq_\alpha C_2$ .

(b) Similar to (a).

(c) If  $\beta_1 < \beta_2$  then any  $\alpha$ -master code  $C_1$  for  $\beta_1$  is  $\beta_2$ -finite; therefore  $\{x \in L_\alpha \mid x \subseteq C_1\}$  and  $\{x \in L_\alpha \mid x \cap C_1 = \emptyset\}$  are  $\beta_2$ -finite, hence  $\Delta_1$  over  $S_{\beta_2}$  and therefore  $\Delta_1$  over  $\langle L_\alpha, C \rangle$  for any  $\alpha$ -master code  $C$  for  $\beta_2$ .

(d) As we have noted, if  $\rho_{n-1}^\beta = \alpha$  then any  $\Sigma_{n-1}$  master code for  $\beta$  is also a  $\Delta_n$  master code for  $\beta$ . The result now follows from (b).

(e) We have  $C \leq_\alpha \alpha\text{-jump}(D)$  as  $C$  is  $\Sigma_1$  over  $\langle L_\alpha, D \rangle$  and  $\alpha\text{-jump}(D)$  has the largest  $\alpha$ -degree of any set  $\Sigma_1$  over  $\langle L_\alpha, D \rangle$ . For the converse ( $\alpha\text{-jump}(D) \leq_\alpha C$ ), it suffices to show that both of the sets  $\{x \in L_\alpha \mid x \subseteq \alpha\text{-jump}(D)\}$  and  $\{x \in L_\alpha \mid x \cap \alpha\text{-jump}(D) = \emptyset\}$  are  $\Sigma_{n+1}$  over  $S_\beta$ . The latter set is actually  $\Pi_1$  over  $\langle L_\alpha, D \rangle$ , hence  $\Pi_n$  over  $S_\beta$ . To complete the proof it suffices to show that  $\{x \in L_\alpha \mid x \subseteq \alpha\text{-jump}(D)\}$  is  $\Sigma_2$  over  $\langle L_\alpha, D \rangle$ .

CLAIM. If  $B \subseteq \alpha$ ,  $B$  is  $\Sigma_1$  over  $\langle L_\alpha, D \rangle$ , then  $B^* = \{x \in L_\alpha \mid x \subseteq B\}$  is  $\Sigma_2$  over  $\langle L_\alpha, D \rangle$ .

*Proof of Claim.* If  $\langle L_\alpha, D \rangle$  is admissible then  $B^*$  is actually  $\Sigma_1$  over  $\langle L_\alpha, D \rangle$ . Otherwise choose  $\gamma < \alpha$ ,  $f: \gamma \rightarrow \alpha$  unbounded such that  $f$  is  $\Sigma_1$  over  $\langle L_\alpha, D \rangle$ . Let  $\varphi(x)$  be a  $\Sigma_1$  formula such that  $y \in B \leftrightarrow \langle L_\alpha, D \rangle \models \varphi(y)$ . Then

$$x \subseteq B \leftrightarrow \exists \text{ sequence } \langle x_\beta \mid \beta < \gamma \rangle \in L_\alpha \text{ such that } x = \bigcup_\beta x_\beta \text{ and} \\ \forall \beta < \gamma \forall y \in x_\beta \langle L_{f(\beta)}, D \cap f(\beta) \rangle \models \varphi(y).$$

This is true since any bounded subset of  $\alpha$  is a member of  $L_\alpha$ . This gives a  $\Sigma_2$  over  $\langle L_\alpha, D \rangle$  definition of  $B^*$ .

(f) Follows from (d) and (e).  $\square$

Our method of showing that every  $\alpha$ -degree  $\geq_\alpha 0'$  is one of the  $d_\nu$ 's,  $\nu \geq 2$  is to prove for each  $\nu$  that for all  $d \geq 0'$ :

$$(*)_\nu \quad d > d_{\nu'} \quad \text{for all } \nu' < \nu \rightarrow d \geq d_\nu.$$

Given this, argue that for all  $d \geq 0'$  there is  $\nu$  such that  $d = d_\nu$  as follows: Let  $\nu$  be least so that  $d \not\geq d_\nu$ . If  $d \neq d_\nu$ , for all  $\nu' < \nu$  then  $d > d_{\nu'}$  for all  $\nu' < \nu$  so by  $(*)_\nu$ ,  $d \geq d_\nu$ , a contradiction.

The proof of  $(*)_\nu$  breaks into cases. The case  $\nu = 2$  is trivial. The successor case is handled by Theorem 4 and Proposition 8 as for any  $\nu$ ,  $d_{\nu+1} = \alpha\text{-jump}(d_\nu)$ . Our next result handles the limit case.

**THEOREM 9.** Suppose  $\rho_{n-1}^\beta > \rho_n^\beta = \alpha$  and let  $C$  be a  $\Sigma_n$  master code for  $\beta$ .

(a) If  $\Sigma_n^p \text{cf}(\rho_{n-1}^\beta) = \omega_1$ , then  $\delta_n^\beta = \alpha$ , there is a  $\Delta_n$  master code  $D$  for  $\beta$  and for all  $A \subseteq \alpha$ :

$$A \text{ is } \beta\text{-finite or } D \leq_\alpha A \vee 0'.$$

(b) If  $\Sigma_n^p \text{cf}(\rho_{n-1}^\beta) \neq \omega_1$  then for all  $A \subseteq \alpha$ :

$$A \text{ is } \beta\text{-finite or } C \leq_\alpha A \vee 0'.$$

For then suppose  $\lambda$  is a limit ordinal,  $\lambda = \omega \cdot \nu$ . Then if (a) above holds with  $\beta = \beta_\nu$ ,  $n = n(\beta_\nu)$ , we have  $d_\lambda = \alpha$ -degree of a  $\Delta_n$  master code for  $\beta_\nu$ ; if (b) holds then  $d_\lambda = \alpha$ -degree of a  $\Sigma_n$  master code for  $\beta_\nu$ . In either case for all  $A \subseteq \alpha$ ,  $A$  is  $\beta$ -finite or  $d_\lambda \leq_\alpha A \vee 0'$ . But  $A$   $\beta_\nu$ -finite implies  $A$  is  $\alpha$ -recursive in some  $\alpha$ -master code for some  $\beta' < \beta_\nu$  and hence  $A \leq_\alpha d_\nu$ , for some  $\nu' < \lambda$ . This demonstrates  $(*)_\lambda$ .

The proof of Theorem 9 necessitates the introduction of some technical notions associated with Skolem functions. Let  $\mathfrak{A} = \langle S_\beta, A \rangle$  be an amenable structure (i.e.,  $A \cap S_\gamma \in S_\beta$  for all  $\gamma < \beta$ ). If  $p \in S_\beta$  then a  $\Sigma_1^p$  Skolem function for  $\mathfrak{A}$  is a partial function  $h(i, x)$  which is  $\Sigma_1$  over  $\mathfrak{A}$  with parameter  $p$  with the property that whenever  $\varphi(x, y, z)$  is  $\Sigma_1(\mathfrak{A})$  then

$$\mathfrak{A} \models \forall x (\exists y \varphi(x, y, p) \longrightarrow \exists i \varphi(x, h(i, x), p)).$$

If  $p = 0$  then we say that  $h$  is a  $\Sigma_1$  Skolem function for  $\mathfrak{A}$ .  $\Sigma_1^p$  Skolem functions are easily constructed (see [0], p. 88). Also we define  $\mathfrak{A}^* = (\beta, A)^* =$  least  $\gamma$  such that there is a  $\Sigma_1(\mathfrak{A})$  function  $f: \beta \xrightarrow{1-1} \gamma$  and  $p(\mathfrak{A}) = p(\beta, A) =$  least  $p$  such that there is such an  $f$  which is  $\Sigma_1$  over  $\mathfrak{A}$  with parameter  $p$ .

We are interested in using a  $\Sigma_1^{p(\beta, A)}$  Skolem function  $h$  to take Skolem hulls and then to analyze the nature of these Skolem hulls after they are transitively collapsed. If  $\gamma < (\beta, A)^*$  then let  $H_\gamma = h[\omega \times \gamma]$  and  $\pi: \langle H_\gamma, \varepsilon \rangle \simeq \langle S_\gamma, \varepsilon \rangle$ . Then  $\gamma$  is  $(\beta, A)$ -pseudostable (or  $\mathfrak{A}$ -pseudostable) if  $\gamma \in H_\gamma$ . It can be easily checked that  $\gamma$   $(\beta, A)$ -pseudostable implies  $\gamma$  is a  $\gamma'$ -cardinal. We also let  $\gamma_A =$  least  $\delta$  such that  $\pi[A \cap H_\gamma]$  is definable over  $S_\delta$ .

In many contexts it is desirable to have a bound on  $\gamma_A$  (in terms of  $\gamma$ ). The reason for this is that if  $\gamma$  is  $(\beta, A)$ -pseudostable and  $B \subseteq (\beta, A)^*$  is  $\Sigma_1$  over  $\mathfrak{A}$  in some parameter in  $\gamma \cup \{p(\beta, A)\}$ , then  $B \cap \gamma$  is  $\Sigma_1$ -definable over  $\langle S_\gamma, \pi[A \cap H_\gamma] \rangle$  and thus a bound on  $\gamma_A$  gives a bound on where  $B \cap \gamma$  is constructed. In case  $(\beta, A)^* = \alpha$  this allows us to estimate the growth rate of  $f_B$ , which is useful in view of Lemma 2.

We describe now an important possible bound on  $\gamma_A$ . For  $\gamma < (\beta, A)^*$ , let  $\hat{\gamma} =$  greatest  $\delta$  such that  $S_\delta \models \gamma$  is a cardinal ( $\hat{\gamma} = \gamma$  if there is no such  $\delta$ ). Then  $A$  is collapsible on  $\beta$  (or  $\mathfrak{A}$  is collapsible) if  $\gamma_A \leq \hat{\gamma}$  for all sufficiently large  $(\beta, A)$ -pseudostable  $\gamma < (\beta, A)^*$ .

*Examples.* (a) If  $A$  is  $\Sigma_1$  over  $S_\beta$  then  $A$  is collapsible. For, choose  $\gamma_0 < (\beta, A)^*$  such that  $A$  is  $\Sigma_1$  over  $S_\beta$  with parameter  $q \in h[\omega \times \gamma_0]$ . Then  $(\gamma \geq \gamma_0, \gamma \text{ is } \beta, A\text{-pseudostable}) \rightarrow A \cap h[\omega \times \gamma]$  is definable over  $h[\omega \times \gamma] \rightarrow \gamma' \geq \gamma_A$ . But  $S_{\gamma'} \models \gamma$  is a cardinal.

(b) If  $A$  is a  $\Delta_n$  ( $\Sigma_{n-1}$ ) master code for  $\mu$  and  $\beta = \delta_n^\mu$  ( $\beta = \rho_{n-1}^\mu$ ) then  $A$  is a collapsible predicate on  $\beta$ .

*Proof.* Choose  $q \in S_\mu$  so that both  $A, \beta - A$  are  $\Sigma_n$  over  $S_\mu$  with parameter  $q$ . Let  $k$  be  $\Sigma_n^q$  Skolem function for  $S_\mu$  and choose  $\gamma_0 < \beta$  so that  $k \cap (\omega \times \beta) \times \beta$  is  $\Sigma_1$  over  $\langle S_\beta, A \rangle$  with some parameter  $r \in S_{\gamma_0}$ . Then for  $\gamma \geq \gamma_0$ ,  $h[\omega \times \gamma]$  is closed under  $k \cap (\omega \times \beta) \times \beta$  and so  $k[\omega \times \gamma] \cap \beta = h[\omega \times \gamma] = H_\gamma$ . Now if  $\gamma$  is  $\beta, A$ -pseudostable let  $K_\gamma = k[\omega \times \gamma]$  and  $\tau: \langle K_\gamma, \varepsilon \rangle \cong \langle S_{\gamma'}, \varepsilon \rangle$ . Then  $\tau \supseteq \pi: \langle H_\gamma, \varepsilon \rangle \cong \langle S_{\gamma'}, \varepsilon \rangle$ ,  $\gamma'' \geq \gamma'$  and  $S_{\gamma''} \models \gamma$  is a cardinal. But  $A \cap K_\gamma$  is  $\Delta_n$  over  $K_\gamma$  so  $\pi[A \cap H_\gamma]$  is  $\Delta_n$  over  $S_{\gamma''}$ .  $\square$

*Proof of Theorem 9.* Let  $\mathfrak{A} = \langle S_{\rho_{n-1}^\beta}, C_{n-1}^\beta \rangle$  and  $\gamma_0 = \Sigma_1 \text{cf}(\mathfrak{A})$ . Choose an  $\mathfrak{A}$ -recursive order-preserving function  $f$  from  $\gamma_0$  onto an unbounded subset of  $\rho_{n-1}^\beta$  such that  $\gamma < \gamma_0, \gamma \text{ limit} \rightarrow \mathfrak{A}_\gamma = \langle S_{f(\gamma)}, C_{n-1}^\beta \cap f(\gamma) \rangle$  is amenable. Let  $s_\delta^i$  be the first  $\mathfrak{A}_\gamma$ -pseudostable greater than  $\aleph_\delta$  (for all  $\gamma < \gamma_0$ ) and let  $s_\delta$  be the first  $\mathfrak{A}$ -pseudostable greater than  $\aleph_\delta$ . Then  $s_\delta = \bigcup \{s_\delta^i \mid \gamma < \gamma_0\}$ .

First assume that  $\gamma_0 < \alpha$ .

Recall that the proof of Lemma 2 shows: If  $\{\delta \mid f_A(\delta) \leq g(\delta)\}$  is stationary and  $g \leq_\alpha B$  then  $A \leq_\alpha B$ . As  $\mathfrak{A}$  is collapsible, whenever  $C \subseteq \alpha$  is  $\Sigma_1$  over  $\mathfrak{A}$  then  $C \cap \aleph_\delta$  is definable over  $S_{s_\delta}^i$  for  $\delta$  sufficiently large. Thus if  $A \subseteq \alpha$  and  $f_A(\delta) \geq s_\delta$  for stationary many  $\delta$  then  $C_n^\beta \leq_\alpha A$ . Otherwise  $f_A(\delta) < s_\delta$  for stationary many  $\delta$  and  $A$  is  $\mathfrak{A}$ -recursive.

(a) Assume  $\gamma_0 = \omega_1$ . Let  $D = \langle s_\delta^i \mid \delta < \omega_1 \rangle$ .

Then  $A \Sigma_1$  over  $\langle S_\alpha, D \rangle \rightarrow A \Sigma_1$  over  $\mathfrak{A}$ . Conversely: By Jensen's extension of embeddings lemma ([1], page 100)  $C_{n-1}^\beta \cap f(\gamma)$  is a  $\Sigma_{n-1}$  master code for some ordinal  $\eta$  such that  $\rho_{n-1}^\beta = f(\gamma)$ , and hence is a collapsible predicate on  $f(\gamma)$ . It follows that if  $\phi(x)$  is  $\Sigma_1$  then  $\{x < \aleph_\delta \mid \mathfrak{A}_\delta \models \phi(x)\}$  is definable over  $S_{\eta_\delta}$  where  $\eta_\delta = \hat{s}_\delta^i$ . Therefore the equivalence:

$$\mathfrak{A} \models \phi(x) \leftrightarrow \exists \delta [x < \aleph_\delta \wedge \mathfrak{A}_\delta \models \phi(x)]$$

shows that any subset of  $\alpha$  which is  $\Sigma_1$  over  $\mathfrak{A}$  is  $\Sigma_1$  over  $\langle S_\alpha, D \rangle$ . So  $D$  is a  $\Delta_1$  master code for  $\mathfrak{A}$  and hence a  $\Delta_n$  master code for  $\beta$ .

Now if  $A \subseteq \alpha$  and  $f_A(\delta) < s_\delta$  for stationary many  $\delta$  then either  $f_A(\delta) < s_\delta^i$  for stationary many  $\delta$  or  $f_A(\delta) \geq s_\delta^i$  for stationary many  $\delta$ . In the former case Fodor's Theorem implies that for some fixed  $\delta_0 < \omega_1$ ,  $f_A(\delta) < s_{\delta_0}^{i_0}$  for stationary many  $\delta$  so  $A$  is  $\rho_{n-1}^\beta$ -finite since the sequence  $\langle s_\delta^{i_0} \mid \delta < \omega_1 \rangle$  is. In the latter case  $D \leq_\alpha A \vee 0'$ . Case (a) is complete.

(b) If  $\gamma_0 \neq \omega_1$ , then  $f_A(\delta) < s_\delta$  for stationary many  $\delta$  implies that there is a fixed  $\delta_0 < \omega_1$  such that  $f_A(\delta) < s_{\delta_0}^{\delta_0}$  for stationary many  $\delta$ . Again  $A$  is  $\rho_{\omega_1}^{\delta_0}$ -finite.

To complete the proof of Theorem 9 we must treat the case  $\gamma_0 > \alpha$ . Then it is still true that  $f_A(\delta) \geq s_\delta$  for stationary many  $\delta$  implies  $C_n^\beta \leq_\alpha A \vee 0'$ . So assume that  $f_A(\delta) < s_\delta$  for stationary many  $\delta$ . Let  $p = p^{(\rho_{\omega_1}^{\delta_0}, C_{n-1}^{\delta_0})}$  and let  $h$  be a  $\Sigma_1^p$  Skolem function for  $\mathfrak{A}$ . For each  $\delta$  such that  $f_A(\delta) < s_\delta$  choose  $g(\delta) < \delta$  so that  $f_A(\delta) \in h[w \times \aleph_{g(\delta)}]$  and  $n_\delta < \omega$ ,  $y_\delta < \aleph_{g(\delta)}$  so that if  $x_\delta = (n_\delta, y_\delta)$  then  $h(x_\delta) = f_A(\delta)$ . Then by Fodor's Theorem there is  $\delta_0 < \omega_1$  such that  $X = \{\delta \mid g(\delta) < \delta_0\}$  is stationary. Then  $\{x_\delta \mid \delta \in X\}$  is  $\alpha$ -finite and as  $\gamma_0 > \alpha$ ,  $h \upharpoonright \{x_\delta \mid \delta \in X\}$  and hence  $A$  is  $\rho_{\omega_1}^{\delta_0}$ -finite.  $\square$

Thus we have established:

**THEOREM.** *For any  $A \subseteq \alpha$ ,  $0' \leq_\alpha A$  has the same  $\alpha$ -degree as an  $\alpha$ -master code.*

### 3. The structure of the $\beta$ -degrees

We continue to assume  $V = L$  and let  $\alpha$  denote  $\aleph_{\omega_1}$ . Now also let  $\beta$  denote a limit ordinal such that:

- (i)  $\beta > \alpha$ , (ii)  $\beta^* = \alpha$ , (iii)  $\Sigma_1 c f \beta < \beta^*$ .

Typical such  $\beta$ 's are  $\alpha \cdot \omega$ ,  $\alpha \cdot \omega_1$ ,  $\alpha \cdot \omega_2$ .

We develop a method of deducing structural properties of the  $\beta$ -degrees from the results of Section 2. If  $e, f$  are  $\beta$ -degrees, define  $e \leq_{w\beta} f$  if  $E \leq_{w\beta} f$  for some  $E \in e$ . Our main result is:

**THEOREM 10.** *There is a well-ordered sequence  $e_0 <_\beta e_1 <_\beta \dots$  of  $\beta$ -degrees of order type  $\aleph_{\omega_1+1}$  such that:*

- (i) For all  $\gamma$ ,  $e_{\gamma+1} \leq_{w\beta} e_\gamma$ .
- (ii) If  $e$  is an arbitrary  $\beta$ -degree then there is a unique  $\gamma$  such that  $e_\gamma \leq_\beta e <_\beta e_{\gamma+1}$ . Also  $E \leq_{w\beta} e_\gamma$  for every  $E \in e$ .

Thus the  $\beta$ -degrees are "nearly" well-ordered. It follows from the second part of (ii) that there do not exist subsets of  $\beta$  which are incomparable with respect to  $\leq_{w\beta}$ .

Our proof of Theorem 10 depends upon choosing special representatives of the  $\alpha$ -degrees. In case  $\Sigma_1 c f \beta = \omega_1$ , all the degrees of Theorem 10 arise as the  $\beta$ -degrees of specially chosen subsets of  $\alpha$ . Moreover this method is of use to us in the case  $\Sigma_1 c f \beta \neq \omega_1$ . For now we only assume that  $\beta$  satisfies conditions (i), (ii), (iii) listed above.

*Definition.* For  $B, C \subseteq \beta$  we say that  $B \leq_{f\beta} C$  if for some  $\beta$ -r.e. sets

$W_0, W_1$ :

$$\begin{aligned} x \subseteq B &\leftrightarrow \exists z, w [\langle x, z, w \rangle \in W_0 \wedge z \subseteq C \wedge w \subseteq \beta - C], \\ y \subseteq \beta - B &\leftrightarrow \exists z, w [\langle y, z, w \rangle \in W_1 \wedge z \subseteq C \wedge w \subseteq \beta - C] \end{aligned}$$

where  $x, y, z, w$  vary over finite subsets of  $\beta$ .

*Definition.* If  $A \subseteq \alpha$  then

$$N_\beta(A) = \{(x, y) \mid x, y \in S_\beta, x \subseteq A, y \subseteq \alpha - A\}$$

and

$$N_\alpha(A) = \{(x, y) \mid x, y \in S_\alpha, x \subseteq A, y \subseteq \alpha - A\}.$$

**LEMMA 11.** *For every  $A \subseteq \alpha$  there is  $A^* \subseteq \alpha$  such that  $A = {}_\alpha A^*$  and  $N_\beta(A^*) \leq_{\beta f} N_\alpha(A^*)$ .*

*Proof.* Let  $p \in S_\beta$  be a parameter such that there are  $f: \beta \xrightarrow{1-1} \alpha$  and  $g: \Sigma_1 cf\beta \rightarrow \beta$  unbounded which are  $\Sigma_1$  over  $S_\beta$  with parameter  $p$ . Let  $\gamma_0 = \Sigma_1 cf\beta$  and choose a  $\Sigma_1^p$  Skolem function  $h(i, x)$  for  $S_\beta$  and approximation  $h^r(i, x)$  such that for  $\gamma < \gamma_0$ ,  $h^r$  is a  $\Sigma_1^p$  Skolem function for  $S_{g(\gamma)}$ . Finally define  $s_\delta^r = h^r[\omega \times \aleph_\delta] \cap \aleph_{\delta+1}$  and  $s_\delta = \bigcup_{r < \gamma_0} s_\delta^r =$  the first  $\beta$ ,  $\phi$ -pseudostable greater than  $\aleph_\delta$ . Note that  $h[\omega \times \alpha] = S_\beta$ .

*Key Fact.* If  $x \in h^r[\omega \times \aleph_\delta]$  and  $x \subseteq \alpha$  then  $f_x(\delta') < s_{\delta'}^{r+1}$  for all  $\delta' \geq \delta$ , where  $f_x =$  cutoff function for  $x$ . (Proof:  $x \cap \aleph_{\delta'} \in h^r[\omega \times \aleph_{\delta'}]$  and so  $f_x(\delta') \leq s_{\delta'}^r < s_{\delta'}^{r+1}$ .)

We assume that  $A$  is not  $\beta$ -finite.

*Case 1.*  $\Sigma_1 cf\beta = \gamma_0 = \omega_1$ . Then the proof of Theorem 9 shows that  $X = \{\delta \mid f_A(\delta) \geq s_\delta^1\}$  is stationary. Define:

$$A^* = \{\gamma \mid \aleph_\delta \leq \gamma \leq f_A(\delta) \text{ for some } \delta \in X\}.$$

*Case 2.*  $\Sigma_1 cf\beta = \gamma_0 \neq \omega_1$ . Then the proof of Theorem 9 shows that  $X = \{\delta \mid f_A(\delta) \geq s_\delta\}$  is stationary. Define:

$$A^* = \{\gamma \mid \aleph_\delta \leq \gamma \leq f_A(\delta) \text{ some } \delta \in X\}.$$

We show now that  $A^*$  works. Clearly  $A^* = {}_\alpha f_A = {}_\alpha A$ . Note that for  $y \subseteq \alpha$ ,  $y \in h^r[\omega \times \aleph_\delta]$ :

$$y \cap [\aleph_\delta, \aleph_{\delta+1}) \neq \emptyset \longrightarrow y \cap [\aleph_\delta, s_\delta^r) \neq \emptyset.$$

Therefore  $y \subseteq \alpha - A^* \rightarrow y = y_1 \cup y_2$  where  $y_1, y_2$  are  $\beta$ -finite and  $y_1 \subseteq \bigcup_{\delta \in X} [\aleph_\delta, \aleph_{\delta+1})$ ,  $y_2$  bounded in  $\alpha$ . Also for  $x \subseteq \alpha$ ,  $x \in h^r[\omega \times \aleph_\delta]$ :

$$x \text{ bounded in } [\aleph_\delta, \aleph_{\delta+1}) \longrightarrow x \subseteq [\aleph_\delta, s_\delta^r).$$

Therefore  $x \subseteq A^* \rightarrow x = x_1 \cup x_2$  where  $x_1, x_2$  are  $\beta$ -finite, for some  $\gamma \in h^r[\omega \times \alpha]$  and

$$x_1 \subseteq \bigcup_{\delta \in X} [\mathfrak{N}_\delta, s_\delta^r], \quad x_2 \text{ bounded in } \alpha.$$

Thus  $N_\beta(A^*) \leq_{f\beta} N_\alpha(A^*)$ . □

Note that for  $A, B \subseteq \alpha$ ,  $A \leq_\alpha B \rightarrow A^* \leq_\beta B$ .

We now apply Lemma 11 to the master code degrees  $d_1, d_2, \dots, d_\gamma, \dots$  which we defined in Section 2. Let  $d_{\gamma_1}$  be the  $\alpha$ -degree of a  $\Sigma_1$  master code for  $\beta$ . Choose canonical representatives  $D_{\gamma_1}, D_{\gamma_1+1}, \dots$  of  $d_{\gamma_1}, d_{\gamma_1+1}, \dots$  respectively and define:

$$\begin{aligned} \hat{e}_0 &= 0, \\ \hat{e}_{1+\gamma} &= \beta\text{-degree } (D_{\gamma_1+1}^*). \end{aligned}$$

Also let  $E_0 = \emptyset$  and  $E_{1+\gamma} = D_{\gamma_1+1}^*$ . Then  $E_\gamma \in \hat{e}_\gamma$  for all  $\gamma$  and  $E_{\gamma+1} =_\alpha \alpha\text{-jump}(E_\gamma)$  for  $\gamma \geq 1$ .

To understand the relationship between  $\hat{e}_\gamma$  and  $\hat{e}_{\gamma+1}$  we must discuss the weak  $\beta$ -jump.

*Definition.* Let  $e, f$  be  $\beta$ -degrees. Then  $f$  is the *weak  $\beta$ -jump* of  $e$  if  $f$  is the largest  $\beta$ -degree which is  $\leq_{w\beta} e$ .

LEMMA 12. (a)  $f$  is the weak  $\beta$ -jump of  $e$  if  $f \leq_\beta g$  for some  $g \leq_{w\beta} e$  and  $e' = \beta\text{-jump}(e) \leq_{w\beta} f$ .

(b)  $\hat{e}_{\gamma+1} =$  the weak  $\beta$ -jump of  $\hat{e}_\gamma$  for all  $\gamma$ .

*Proof.* (a) Choose representatives  $E, F$  of  $e, f$ , respectively. It suffices to show that  $E' \leq_{w\beta} F, G \leq_{w\beta} E \rightarrow G \leq_\beta F$ . Choose  $e_1, e_2$  so that

$$\begin{aligned} x \in G &\leftrightarrow \{e_1\}_\beta^E(x) \text{ diverges,} \\ x \notin G &\leftrightarrow \{e_2\}_\beta^E(x) \text{ diverges.} \end{aligned}$$

Also choose a  $\beta$ -recursive  $f(K, e)$  such that  $\{f(K, e)\}_\beta^E(0)$  diverges  $\leftrightarrow \forall x \in K, \{e\}_\beta^E(x)$  diverges. Then

$$\begin{aligned} K \subseteq G &\leftrightarrow \{f(K, e_1)\}_\beta^E(0) \text{ diverges,} \\ H \subseteq \beta - G &\leftrightarrow \{f(H, e_2)\}_\beta^E(0) \text{ diverges,} \end{aligned}$$

so  $\beta$ -finite neighborhood questions about  $G$  can be answered using finite neighborhood information on  $E'$ . But  $E' \leq_{w\beta} F$  so  $G \leq_\beta F$ .

(b) Recall that  $N_\beta(E_\gamma) \leq_{f\beta} N_\alpha(E_\gamma)$ . We claim that  $\beta\text{-jump}(E_\gamma) \leq_{w\beta} E_{\gamma+1}$ . In case  $\gamma = 0$  this is clear as  $E_1$  is a  $\Sigma_1$  master code for  $\beta$ . Otherwise note that any subset of  $\alpha$  which is  $\Sigma_1$  over  $S_\beta$  is  $\Sigma_1$  over  $\langle S_\alpha, E_\gamma \rangle$  so we can write:

$$\begin{aligned} (e, x) \in \beta\text{-jump}(E_\gamma) &\leftrightarrow \exists z, w \in S_\beta[\langle x, z, w \rangle \in W_e \wedge z \subseteq E_\gamma \wedge w \subseteq \beta - E_\gamma] \\ &\leftrightarrow \exists z, w \in S_\alpha[\langle x, z, w \rangle \in W_{f(e)} \wedge z \subseteq E_\gamma \wedge w \subseteq \alpha - E_\gamma] \end{aligned}$$

where  $f$  is  $\beta$ -recursive. This last predicate is  $\Sigma_1$  over  $\langle S_\alpha, E_\gamma \rangle$ . Thus  $\beta\text{-jump}(E_\gamma) \leq_{f\beta} \alpha\text{-jump}(E_\gamma) =_\alpha E_{\gamma+1}$ . So  $\beta\text{-jump}(E_\gamma) \leq_{w\beta} E_{\gamma+1}$ .

Now define:  $\hat{E}_\gamma = \{\langle \delta, e, x \rangle \mid \{e\}_\alpha^{E_\gamma}(x) \text{ converges by stage } \aleph_\delta\} \subseteq \omega_1 \times \alpha \times \alpha$ . (Here,  $\{e\}_\alpha^E$  denotes the  $e$ th function partial  $\alpha$ -recursive in  $E$ .) Then  $\hat{E}_\gamma \leq_{w\alpha} E_\gamma$  so  $\hat{E}_\gamma \leq_\alpha \alpha\text{-jump}(E_\gamma)$ . Also  $\alpha\text{-jump}(E_\gamma) \leq_\alpha \hat{E}_\gamma$ . For  $K \in S_\alpha$ ,

$$\begin{aligned} K \subseteq \alpha\text{-jump}(E_\gamma) &\leftrightarrow \forall x \in K, \quad (x = \langle x_0, x_1 \rangle \text{ and } \{x_0\}_\alpha^{E_\gamma}(x_1) \text{ converges}) \\ &\leftrightarrow \exists f \in S_\alpha, \quad (f: \omega_1 \longrightarrow \alpha \wedge \bigcup \text{Range } f = K \wedge \forall x \in f(\delta), \\ &\quad \{x_0\}_\alpha^{E_\gamma}(x_1) \text{ converges by stage } \aleph_\delta, \\ &\quad \text{for every } \delta < \omega_1) \\ &\leftrightarrow \exists f \in S_\alpha, \quad (f: \omega_1 \longrightarrow \alpha \wedge K = \bigcup \text{Range } f \wedge \forall \delta < \omega_1, \\ &\quad \forall x \in f(\delta), \langle \delta, x_0, x_1 \rangle \in \hat{E}_\gamma). \end{aligned}$$

And, for  $H \in S_\alpha$ ,

$$\begin{aligned} H \subseteq \alpha - (\alpha\text{-jump}(E_\gamma)) &\leftrightarrow \forall x \in H, \quad (\{x_0\}_\alpha^{E_\gamma}(x_1) \text{ diverges}) \\ &\leftrightarrow \forall x \in H, \quad (\omega_1 \times \{x_0\} \times \{x_1\} \subseteq \alpha - \hat{E}_\gamma). \end{aligned}$$

So  $\alpha\text{-jump}(E_\gamma) =_\alpha \hat{E}_\gamma$ .

Now we have  $E_{\gamma+1} =_\alpha \alpha\text{-jump}(E_\gamma) =_\alpha \hat{E}_\gamma$  so  $E_{\gamma+1} \leq_\beta \hat{E}_\gamma$  (since for any  $B \subseteq \alpha$ ,  $E_{\gamma+1} \leq_\alpha B \leftrightarrow E_{\gamma+1} \leq_\beta B$ ). But  $\beta\text{-jump}(E_\gamma) \leq_{w\beta} E_{\gamma+1}$  so the hypotheses of part (a) are now satisfied and  $\hat{e}_{\gamma+1} = \beta\text{-degree}(E_{\gamma+1}) = \text{weak } \beta\text{-jump}(\hat{e}_\gamma)$ .  $\square$

For all inadmissible  $\beta$  and all  $\beta$ -degrees  $e$ ,  $\text{weak } \beta\text{-jump}(\text{weak } \beta\text{-jump}(e)) = \beta\text{-jump}(e)$ . Thus in this case it is appropriate to refer to  $\text{weak } \beta\text{-jump}$  as the “ $\beta$ -half-jump.”

We now have all the ingredients needed to provide a proof of Theorem 10 in the case  $\Sigma_1 \text{cf } \beta = \omega_1$ . For such a  $\beta$  let  $e_\gamma = \hat{e}_\gamma$  for each  $\gamma \geq 0$ . As  $\Sigma_1 \text{cf } \beta = \omega_1$  there is a tame  $\Sigma_1(S_\beta)$  bijection  $f: \beta \leftrightarrow \alpha$ ; i.e., for each  $\delta < \alpha$ ,  $f^{-1}[\delta]$  is  $\beta$ -finite and the sequence  $\langle f^{-1}[\delta] \mid \delta < \alpha \rangle$  is  $\Sigma_1$  over  $S_\beta$ .

Now let  $B \subseteq \beta$  and consider  $f[B] \subseteq \alpha$ . Then by the proof of Theorem 9 either  $f[B]$  is  $\beta$ -recursive or  $E_1 \leq_\alpha f[B]$ . In the former case,  $B$  is  $\beta$ -recursive and so  $0 = e_0 \leq \beta\text{-degree}(B) \leq e_1$  and  $B \leq_{w\beta} e_0$ . In the latter case choose  $\gamma$  so that  $E_\gamma =_\alpha f[B]$ . Then  $N_\alpha(f[B]) \leq_{w\beta} B$  as  $f$  is tame and so we have

$$N_\beta(E_\gamma) \leq_{f\beta} N_\alpha(E_\gamma) \leq_{f\beta} N_\alpha(f[B]) \leq_{w\beta} B;$$

therefore  $E_\gamma \leq_\beta B$ . Also  $B \leq_{f\beta} f[B] \leq_{w\beta} E_\gamma$  so  $B \leq_{w\beta} E_\gamma$ . This proves Theorem 10 in this case.

*Note.* There is an easier proof of Theorem 10 in case  $\Sigma_1 \text{cf } \beta = \omega_1$ : Define a  $\beta$ -degree  $d$  to be *regular* if  $\langle S_\beta, D \rangle$  is amenable for some  $D \in d$ . One can show that in case  $\Sigma_1 \text{cf } \beta = \omega_1$ , the regular  $\beta$ -degrees are well-ordered and for any  $A \subseteq \beta$  there is a regular  $\beta$ -degree  $d$  such that  $A \leq_{w\beta} d$  and  $d \leq_\beta \beta\text{-degree}(A)$ . Moreover if  $d$  is regular, then  $\text{weak } \beta\text{-jump}(d)$  is the least regular  $\beta$ -degree greater than  $d$ . However the proof that we have given shows that

the degrees  $e_\gamma$  have representatives contained in  $\alpha$  and also provides us with the objects needed to analyze the  $\beta$ -degrees when  $\Sigma_1 cf\beta \neq \omega_1$ .

Finally we establish Theorem 10 in the case  $\Sigma_1 cf\beta \neq \omega_1$ . Recall the degrees  $\hat{e}_\gamma$ , obtained by taking the  $\beta$ -degrees of specially chosen representatives of the  $\alpha$ -degrees  $\geq \alpha$ -degree ( $\Sigma_1$  master code for  $\beta$ ). In our present case ( $\Sigma_1 cf\beta \neq \omega_1$ ), the conclusion of Theorem 10 does not hold if we simply take  $e_\gamma = \hat{e}_\gamma$ ; other natural  $\beta$ -degrees arise. For any limit ordinal  $\gamma$  define  $\lambda$  so that  $E_\gamma \in \hat{e}_\gamma$  is an  $\alpha$ -master code for  $\lambda$  and let  $\mathfrak{A}_\gamma = \langle S_{\rho_{n-1}^\lambda}, C_{n-1}^\lambda \rangle$  where  $n = n(\lambda) =$  least  $m$  such that  $\rho_m^\lambda$  equals  $\alpha$ ,  $C_{n-1}^\lambda = \Sigma_{n-1}$  master code for  $\lambda$ . Then  $\gamma$  is *masterful* if  $\Sigma_1 cf \mathfrak{A}_\gamma = \gamma_0 = \Sigma_1 cf\beta$ . We shall define a certain  $\beta$ -degree  $\hat{f}_\gamma$  for masterful  $\gamma$ .

LEMMA 13.  $\gamma$  masterful  $\rightarrow \Sigma_1^{\aleph_\gamma}$ -cofinality( $\gamma$ ) =  $\gamma_0$ .

*Proof.* Note that  $\gamma = \omega \cdot \text{ordertype} \{ \beta' < \rho_{n-1}^\lambda \mid \beta' \text{ is projectible into } \alpha \}$ .

(a) If  $\alpha =$  largest  $\rho_{n-1}^\lambda$ -cardinal then  $P = \{ \beta' < \rho_{n-1}^\lambda \mid \beta' \text{ is projectible into } \alpha \}$  is unbounded in  $\rho_{n-1}^\lambda$  unless  $\rho_{n-1}^\lambda$  is of the form  $\beta' + \omega$ . In the former case

$$\begin{aligned} \Sigma_1^{\aleph_\gamma}\text{-cofinality}(\gamma) &= \Sigma_1^{\aleph_\gamma}\text{-cofinality}(\text{ordertype of } P) \\ &= \Sigma_1^{\aleph_\gamma}\text{-cofinality}(\rho_{n-1}^\lambda) = \gamma_0 \end{aligned}$$

and in the latter case

$$\Sigma_1^{\aleph_\gamma}\text{-cofinality}(\gamma) = \omega = \Sigma_1^{\aleph_\gamma}\text{-cofinality}(\rho_{n-1}^\lambda) = \gamma_0 .$$

(b) If  $\kappa =$  next  $\rho_{n-1}^\lambda$ -cardinal after  $\alpha$  then  $\gamma = \kappa$ . There exists a parameter  $p \in S_{\rho_{n-1}^\lambda}$  such that if  $h$  is a  $\Sigma_1^p$  Skolem function for  $\mathfrak{A}_\gamma$  then  $h[\omega \times \alpha] = S_{\rho_{n-1}^\lambda}$ . Let  $f: \gamma_0 \rightarrow \rho_{n-1}^\lambda$  be unbounded and  $\Sigma_1$  over  $\mathfrak{A}_\gamma$ . Then define  $g: \gamma_0 \rightarrow \kappa$  by  $g(\delta) = \sup \{ h^{f(\delta)}[\omega \times \alpha] \cap \kappa \}$ , where  $h^{f(\delta)}$  is the interpretation of a  $\Sigma_1^p$  definition of  $h$  inside  $\langle S_{f(\delta)}, C_{n-1} \cap f(\delta) \rangle$  (if this structure is amenable then  $h^{f(\delta)}$  is a  $\Sigma_1^p$  Skolem function for it). The function  $g$  is  $\Sigma_1$  over  $\mathfrak{A}_\gamma$  and unbounded since  $\bigcup_\delta h^{f(\delta)} = h$  and  $\kappa \subseteq h[\omega \times \alpha]$ .  $\square$

Now let  $f_\gamma: \gamma_0 \rightarrow \gamma$  be cofinal, continuous, increasing and  $\Sigma_1(\mathfrak{A}_\gamma)$ . Also choose  $f: \gamma_0 \rightarrow \beta$  cofinal, continuous, increasing and  $\Sigma_1(S_\beta)$ . Define:

$$F_\gamma = \{ \langle f(\gamma'), \delta \rangle \mid \gamma' < \gamma_0 \wedge \delta \in E_{f_\gamma(\gamma')} \} \subseteq S_\beta .$$

(Thus  $F_\gamma$  is obtained by "spreading out" the sequence  $\langle E_{f_\gamma(\gamma')} \mid \gamma' < \gamma_0 \rangle$  cofinally in  $S_\beta$ .)  $F_\gamma$  is the  $\Sigma_1(S_\beta)$  union of:

$$F_\gamma^{\gamma'} = \{ \langle f(\gamma'), \delta \rangle \mid \delta \in E_{f_\gamma(\gamma')} \} , \quad \gamma' < \gamma_0$$

and for  $\beta' < \beta$ ,  $S_{\beta'} \cap F_\gamma^{\gamma'} = \emptyset$  for sufficiently large  $\gamma' < \gamma_0$ .

LEMMA 14. For all masterful  $\gamma$ ,



$$\begin{aligned} N_{\beta}(F_{\gamma}) &= \{(x, y) \mid x, y \in S_{\beta} \wedge x \subseteq F_{\gamma} \wedge y \subseteq S_{\beta} - F_{\gamma}\} \\ &\leq_{w\beta} N_{\beta}(F_{\gamma}) \cap \{(x, y) \mid x \text{ and } y \text{ are contained in } f[\gamma_0] \times \delta, \\ &\quad \text{for some } \delta < \alpha\}. \end{aligned}$$

*Proof.* We employ the Skolem function  $h$  introduced in the proof of Lemma 11. Then for all  $\gamma' < \gamma_0$  there is a stationary set  $X_{\gamma'} \subseteq \omega_1$  such that if  $y \subseteq \alpha - E_{f_{\gamma'}(\gamma')}$ ,  $y \in h[\omega \times \aleph_{\nu}]$  then  $y = y_1 \cup y_2$  where  $y_1, y_2$  are  $\beta$ -finite,  $y_1 \subseteq \bigcup_{\nu' \in X_{\gamma'}} [\aleph_{\nu'}, \aleph_{\nu'+1})$  and  $y_2 \subseteq \aleph_{\nu}$ . It follows that if  $H \subseteq S_{\beta} - F_{\gamma}$ ,  $H \in h[\omega \times \aleph_{\nu}]$  then  $H = H_1 \cup H_2$  where  $H_1, H_2$  are  $\beta$ -finite,

$$H_1 \subseteq \bigcup_{\gamma' < \gamma_0} (\{f(\gamma')\} \times \bigcup_{\nu' \in X_{\gamma'}} [\aleph_{\nu'}, \aleph_{\nu'+1})) \quad \text{and} \quad H_2 \subseteq f[\gamma_0] \times \aleph_{\nu}.$$

Also if  $x \subseteq \alpha$ ,  $x \in h[\omega \times \aleph_{\nu}]$ , then  $x \subseteq E_{f_{\gamma'}(\gamma')}$  if and only if  $x \subseteq \bigcup_{\nu' \in X_{\gamma'}} [\aleph_{\nu'}, \aleph_{\nu'+1})$ ,  $|x \cap \aleph_{\nu'+1}| \leq \aleph_{\nu'}$  for each  $\nu' < \omega_1$  and  $x \cap \aleph_{\nu} \subseteq E_{f_{\gamma'}(\gamma')}$ . It follows that if  $K \subseteq S_{\beta}$ ,  $K \in h[\omega \times \aleph_{\nu}]$  then  $K \subseteq F_{\gamma}$  if and only if  $K \subseteq \bigcup_{\gamma' < \gamma_0} (\{f(\gamma')\} \times \bigcup_{\nu' \in X_{\gamma'}} [\aleph_{\nu'}, \aleph_{\nu'+1}))$ ,  $|K \cap \{f(\gamma')\} \times \aleph_{\nu'+1}| \leq \aleph_{\nu'}$  for  $\nu' < \omega_1$ ,  $\gamma' < \gamma_0$  and  $K \cap f[\gamma_0] \times \aleph_{\nu} \subseteq F_{\gamma}$ . These two facts suffice to prove the lemma.  $\square$

We are ready to define the  $e_{\gamma}$ 's in this case. For  $\gamma$  masterful let  $\hat{f}_{\gamma} = \beta$ -degree( $F_{\gamma}$ ). Then we set, for  $\gamma = 0$  or a limit,  $n \in \omega$ :

$$\begin{aligned} e_{\gamma+n} &= \hat{e}_{\gamma+n}, & \gamma \text{ not masterful} \\ e_{\gamma} &= \hat{f}_{\gamma}, & e_{\gamma+n+1} = \hat{e}_{\gamma+n} \quad \gamma \text{ masterful}. \end{aligned}$$

Our next lemma is the key lemma toward understanding the degrees  $e_{\gamma}$ ,  $\gamma$  masterful. Let  $g_{\gamma}: \gamma_0 \rightarrow \rho_{n-1}^{\lambda}$  be cofinal, continuous, increasing and  $\Sigma_1(\mathcal{Q}_{\gamma})$ . In case  $\gamma_0 > \omega$  we also assume that  $C_{n-1}^{\lambda} \cap g_{\gamma}(\gamma')$  is a regular subset of  $g_{\gamma}(\gamma')$  for each  $\gamma' < \gamma_0$ . If  $\gamma_0 = \omega$  then there is a  $\Sigma_{n-1}$  master code  $D_{n-1}^{\lambda}$  for  $\lambda$  such that  $g_{\gamma}(\gamma') \cap D_{n-1}^{\lambda}$  is finite (and hence collapsible) for each  $\gamma' < \omega$ . Finally let  $\bar{C} = C_{n-1}^{\lambda}$  if  $\gamma_0 > \omega$ ,  $\bar{C} = D_{n-1}^{\lambda}$  if  $\gamma_0 = \omega$  and set  $s(\nu, \gamma') = \text{least } g_{\gamma}(\gamma')$ ,  $\bar{C} \cap g_{\gamma}(\gamma')$ -pseudostable greater than  $\aleph_{\nu}$ , for each  $\nu < \omega_1$ ,  $\gamma' < \gamma_0$ .

**UNIFORMITY LEMMA.** (a) *If  $s \leq_{w\beta} B$  then  $F_{\gamma} \leq_{\beta} B$ .*

(b) *If  $t: \omega_1 \times \gamma_0 \rightarrow \alpha$ ,  $t \leq_{w\beta} B$  and for each  $\gamma' < \gamma_0$ ,  $\{\nu \mid s(\nu, \gamma') < t(\nu, \gamma')\}$  is stationary then  $s \leq_{w\beta} B$ .*

*Proof.* (a) If  $A \subseteq \alpha$  let  $f_A: \omega_1 \rightarrow \alpha$  denote the cutoff function for  $A$  (as defined in Section 1). For  $\gamma' < \gamma_0$ , let  $G_{\gamma'} = \bigcup_{\gamma'' < \gamma'} \{\gamma''\} \times E_{f_{\gamma'}(\gamma'')} \subseteq \alpha$  and let  $g_{\gamma'} = f_{G_{\gamma'}}$ . Then  $g_{\gamma'}$  is  $\rho_{n-1}^{\lambda}$ -finite so there is an ordinal  $h(\gamma') < \gamma_0$  such that  $g_{\gamma'}(\nu) < s(\nu, h(\gamma'))$  for sufficiently large  $\nu < \omega_1$ . By Lemma 14 it suffices to show that  $g \leq_{w\beta} B$  where  $g(\nu, \gamma') = g_{\gamma'}(\nu)$ . But by the way we have defined  $G_{\gamma'}$ , it actually suffices to show that  $g \upharpoonright X \leq_{w\beta} B$  where for unboundedly many  $\gamma' < \gamma_0$ ,  $\{\nu \mid \langle \nu, \gamma' \rangle \in X\}$  is unbounded.

We repeat now the argument of Lemma 2. For each  $\gamma' < \gamma_0$  and  $\nu < \omega_1$ , let  $m_{\nu}^{\gamma'}$  be the  $<_J$ -least injection of  $s(\nu, h(\gamma'))$  into  $\aleph_{\nu}$ , and choose a stationary set  $X_{\gamma'} \subseteq \omega_1$  such that  $Y_{\gamma'} = \{m_{\nu}^{\gamma'}(g_{\gamma'}(\nu)) \mid \nu \in X_{\gamma'}\}$  is a bounded subset of  $\alpha$ . As  $\gamma_0 \neq \omega_1$ , there is an unbounded  $Y \subseteq \gamma_0$  such that  $\bigcup_{\gamma' \in Y} Y_{\gamma'}$  is bounded in  $\alpha$ . But then  $g \upharpoonright X \leq_{f_\alpha} s$  where  $X = \{\langle \nu, \gamma' \rangle \mid \nu \in X_{\gamma'} \wedge \gamma' \in Y\}$ . Since  $s \leq_{w_\beta} B$  we get  $g \upharpoonright X \leq_{w_\beta} B$ .

(b) Note that for all  $\nu < \omega_1$ ,  $\gamma' < \gamma_0$ ,  $s \upharpoonright \nu \times \gamma' \in S_\eta$  where  $\eta = \text{largest } \eta'$  such that  $S_{\eta'} \models s(\nu, \gamma')$  is a cardinal (use the fact that  $\bar{C} \cap g_{\gamma'}(\gamma')$  is a collapsible predicate on  $g_{\gamma'}(\gamma')$ ). Now define  $\tilde{s}(\nu, \gamma') = \eta$  if  $s \upharpoonright \nu \times \gamma'$  is the  $\eta$ th set in the canonical well-ordering  $<_J$ . Thus we can assume that for each  $\gamma' < \gamma_0$ ,  $\{\nu \mid \tilde{s}(\nu, \gamma') < t(\nu, \gamma')\}$  is stationary and it suffices to produce  $X \subseteq \omega_1 \times \gamma_0$  such that  $\tilde{s} \upharpoonright X \leq_{w_\beta} B$ , and for unboundedly many  $\gamma' < \gamma_0$ ,  $\{\nu \mid \langle \nu, \gamma' \rangle \in X\}$  is unbounded in  $\omega_1$ . Now simply repeat the argument in the second paragraph of part (a) to get such an  $X$ .  $\square$

**COROLLARY TO PROOF.** *Suppose  $t: \omega_1 \times \gamma_0 \rightarrow \alpha$ ,  $u: \omega_1 \times \gamma_0 \rightarrow \alpha$  and for each  $\gamma' < \gamma_0$ ,  $\{\nu \mid t(\nu, \gamma') < u(\nu, \gamma')\}$  is stationary. Then  $t \upharpoonright X \leq_{f_\alpha} u$  for some  $X$  such that for unboundedly many  $\gamma' < \gamma_0$ ,  $\{\nu \mid \langle \nu, \gamma' \rangle \in X\}$  is stationary.*

The Uniformity Lemma and its corollary are key steps in completing our proof. We first establish the relationship between  $\hat{e}_\gamma$  and  $\hat{f}_\gamma$  for masterful  $\gamma$ .

**LEMMA 15.** *If  $\gamma$  is masterful then  $\hat{e}_\gamma = \text{weak } \beta\text{-jump } \hat{f}_\gamma$ .*

*Proof.* We use Lemma 12(a). Let  $E_\gamma$  and  $F_\gamma$  be as defined earlier. If  $k: \beta \leftrightarrow \alpha$  is a  $\beta$ -recursive bijection then  $E_\gamma \leq_\beta k[F_\gamma]$  since  $F_\gamma$  (and hence  $k[F_\gamma]$ ) is not  $\rho_{n-1}^1$ -finite. As  $k[F_\gamma] \leq_{w_\beta} F_\gamma$  it only remains to show  $\beta\text{-jump}(F_\gamma) \leq_{w_\beta} E_\gamma$  in order to establish our lemma.

Note that the sequence  $\langle E_{f_{\gamma'}(\gamma')} \mid \gamma' < \gamma_0 \rangle$  is  $\Sigma_1(\mathfrak{A}_\gamma)$  and therefore so is the sequence  $\langle F_\gamma \cap f(\gamma') \mid \gamma' < \gamma_0 \rangle$ . It follows that  $N_\beta(F_\gamma)$  is  $\Sigma_1$  over  $\mathfrak{A}_\gamma$  and since

$$\langle e, x \rangle \in \beta\text{-jump}(F_\gamma) \leftrightarrow \exists \langle z, w \rangle \in N_\beta(F_\gamma) \quad [\langle x, z, w \rangle \in W_e^\beta]$$

we see that  $\beta\text{-jump}(F_\gamma)$  is  $\Sigma_1$  over  $\mathfrak{A}_\gamma$ . But  $E_\gamma$  is a  $\Sigma_n$  master code for  $\lambda$  (a  $\Sigma_1$  master code for  $\mathfrak{A}_\gamma$ ) and therefore  $\beta\text{-jump}(F_\gamma) \leq_{f_\beta} k[\beta\text{-jump}(F_\gamma)] \leq_\alpha E_\gamma$ . Thus  $\beta\text{-jump}(F_\gamma) \leq_{w_\beta} E_\gamma$ .  $\square$

We now complete our proof. First choose an increasing  $\beta$ -recursive sequence  $\langle K_\gamma \mid \gamma' < \gamma_0 \rangle$  of  $\beta$ -finite sets such that  $\beta$ -cardinality  $(K_\gamma) = \alpha$  for all  $\gamma'$  and  $S_\beta = \bigcup_\gamma K_\gamma$  (using the facts that  $\beta^* = \alpha$  and  $\gamma_0 = \Sigma_1 \text{cf } \beta$ ). For each  $\gamma' < \gamma_0$  let  $k_{\gamma'}$  be the  $<_J$ -least bijection from  $K_\gamma$  onto  $\alpha$ .

Now choose  $B \subseteq S_\beta$  and let  $B_{\gamma'} = k_{\gamma'}[B \cap K_\gamma]$ . Define  $g: \gamma_0 \rightarrow \aleph_{\omega_1+1}$  by the property:  $B_{\gamma'} =_\alpha E_{g(\gamma')}$ . Let  $\gamma = \sup g[\gamma_0]$ .

*Case 1.*  $\gamma = g(\gamma')$  for some  $\gamma'$ . If  $\gamma = 0$  let  $\lambda = \beta$  and  $n = 1$ . Otherwise, let  $E_\gamma$  be a  $\Delta_n$  master code for  $\lambda$ . Let  $\delta = \Sigma_1 \text{cf} \mathfrak{N}_{n-1}^i$  and choose  $s: \omega_1 \times \delta \rightarrow \alpha$  to be  $\Sigma_n$  over  $S_i$  and such that  $\lim_{\delta' \rightarrow \delta} s(\nu, \delta') = \text{least } \rho_{n-1}^i, C_{n-1}^i\text{-pseudostable} > \aleph_i$  for each  $\nu < \omega_1$ . For each  $\gamma' < \gamma_0$ , if  $j_{\gamma'} =$  cutoff function for  $B_{\gamma'}$  then  $\{\nu \mid j_{\gamma'}(\nu) < s(\nu, \delta')\}$  is stationary for some  $\delta' < \delta$ . It follows from the corollary to the Uniformity Lemma that if we let  $j(\nu, \gamma') = j_{\gamma'}(\nu)$  then  $j \upharpoonright X \leq_{w\beta} E_\gamma$  for some  $X \subseteq \omega_1 \times \gamma_0$  such that for unboundedly many  $\gamma' < \gamma_0$ ,  $\{\nu \mid (\nu, \gamma') \in X\}$  is stationary. Thus since  $\langle K_{\gamma'} \mid \gamma' < \gamma_0 \rangle$  is increasing,  $j \leq_{w\beta} E_\gamma$  and so  $B \leq_{w\beta} E_\gamma$ . Also  $E_\gamma \leq_\beta B$  since for some  $\gamma' < \gamma_0$ ,  $E_\gamma \leq_\alpha B_{\gamma'}$  (so  $N_\beta(E_\gamma) \leq_{f\beta} N_\alpha(E_\gamma) \leq_{f\beta} N_\alpha(B_{\gamma'}) \leq_{w\beta} B$ ).

*Case 2.* Otherwise. Choose  $\lambda$  so that  $E_\gamma$  is a master code for  $\lambda$  and let  $n$  be least so that  $\rho_n^i = \alpha$ . Let  $j_{\gamma'} =$  the cutoff function for  $B_{\gamma'}$  and  $j(\nu, \gamma') = j_{\gamma'}(\nu)$ . Also define  $\delta = \Sigma_1 \text{cf}(\mathfrak{N}_{n-1}^i)$  and choose  $s: \omega_1 \times \delta \rightarrow \alpha$  to be  $\Sigma_n(\mathfrak{N}_{n-1}^i)$  and such that  $\lim_{\delta' \rightarrow \delta} s(\nu, \delta') =$  the least  $\rho_{n-1}^i, C_{n-1}^i\text{-pseudostable}$  greater than  $\aleph_i$ . As each  $B_{\gamma'}$  is  $\rho_{n-1}^i$ -finite, for each  $\gamma' < \gamma_0$  there is  $\delta' < \delta_0$  such that  $\{\nu \mid j(\nu, \gamma') < s(\nu, \delta')\}$  is stationary. But it cannot be the case that for a fixed  $\delta' < \delta$ ,  $\{\nu \mid j(\nu, \gamma') < s(\nu, \delta')\}$  is stationary for unboundedly many  $\gamma' < \gamma_0$ . Thus if  $\delta < \alpha$  we must have  $\delta = \gamma_0$ . And, the argument at the end of the proof of Theorem 9 shows that  $\delta < \alpha$ . Thus  $\gamma$  is masterful.

As  $\gamma_0 \neq \omega_1$  there is no  $\Delta_n$  master code for  $\lambda$  and thus  $E_\gamma$  is a  $\Sigma_n$  master code for  $\lambda$ . Now for each  $\gamma' < \gamma_0$  choose  $h_1(\gamma')$  so that  $\{\nu \mid s(\nu, \gamma') < j(\nu, h_1(\gamma'))\}$  contains a closed unbounded set (by the preceding paragraph). The Uniformity Lemma (b) is satisfied by  $t(\nu, \gamma') = j(\nu, h_1(\gamma'))$  and thus  $F_\gamma \leq_\beta B$ . But the above argument applies equally well to  $F_\gamma$  as to  $B$ , so for each  $\gamma' < \gamma_0$  choose  $h_2(\gamma')$  so that  $\{\nu \mid s(\nu, \gamma') < l(\nu, h_2(\gamma'))\}$  contains a closed, unbounded set where  $l(\nu, \gamma') = l_{\gamma'}(\nu)$  and  $l_{\gamma'}$  is the cutoff function for  $k_{\gamma'}[F_\gamma \cap K_{\gamma'}]$ . Then for each  $\gamma' < \gamma_0$  there is  $h_3(\gamma')$  such that  $\{\nu \mid j(\nu, \gamma') < l(\nu, h_3(\gamma'))\}$  is stationary so the corollary to the Uniformity Lemma applies to show that for some “large”  $X$ ,  $j \upharpoonright X \leq_{f\alpha} l$ . But  $j \leq_{f\beta} j \upharpoonright X$  and  $l \leq_{w\beta} F_\gamma$  so  $j \leq_{w\beta} F_\gamma$ . Thus  $B \leq_{w\beta} F_\gamma$ .  $\square$

We have established the conclusion of Theorem 10 when  $\Sigma_1 \text{cf} \beta \neq \omega_1$ .

#### 4. Further results and open questions

As we have earlier mentioned, there are incomparable  $\aleph_\omega$ -degrees above  $0'$  as well as infinite descending sequences of  $\aleph_\omega$ -degrees. However the degrees in these examples are of sets  $\Delta_2$  over  $L_{\aleph_\omega^+}$  where  $\aleph_\omega^+$  = next admissible after  $\aleph_\omega$ . Thus we propose:

*Problem 1.* Show that there are incomparable  $\aleph_\omega$ -degrees above  $0'$  and

infinite descending sequences of  $\aleph_\omega$ -degrees above  $0'$  constructed (in  $L$ ) before  $\aleph_\omega^+$ .

Harrington has shown that if  $\Sigma_1 cf\beta = \Sigma_1^{\beta} cf(\beta^*) < \beta^*$  then incomparable  $\beta$ -r.e. degrees exist. We have extended this to show that for any such  $\beta$  there are incomparable  $\beta$ -degrees between  $e_\gamma$  and  $e_{\gamma+1}$  for all  $\gamma$  (where  $e_\gamma$  is as in Theorem 10). These results will be presented in [3]. However the case  $\Sigma_1 cf\beta \neq \Sigma_1^{\beta} cf(\beta^*)$  remains unsettled.

*Problem 2.* Show that if  $\Sigma_1 cf\beta, \Sigma_1^{\beta} cf(\beta^*) < \beta^*, \Sigma_1 cf\beta \neq \Sigma_1^{\beta} cf(\beta^*)$  then there exist incomparable  $\beta$ -r.e. degrees.

Our last problem concerns the optimality of Theorem 10. A positive solution shows that  $\{e_\gamma \mid \gamma < \aleph_{\omega_1+1}\}$  is definable over the  $\beta$ -degrees as a partial  $\exists$  ordering. For then this collection consists of all  $\beta$ -degrees  $e$  such that for all  $\beta$ -degrees  $f$ , either  $f \leq e$  or  $e \leq f$ .

*Problem 3.* Let  $\langle e_\gamma \mid \gamma < \aleph_{\omega_1+1} \rangle$  be defined as in Theorem 10. Show that if  $e_\gamma < f < e_{\gamma+1}$ , then there is a  $g$  incomparable  $f$ ,  $e_\gamma < g < e_{\gamma+1}$ .

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